

TROTTER CIRCUIT OPTIMIZATION

THROUGH ADIABATIC COMPUTATION

Kim, Hyunseong

Quantum Field and Gravity Theory Group,
GIST, Department of Physics and Photonics

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Part I

QUBO PROBLEMS IN CIRCUIT OPTIMIZATION

LIE-TROTTER FORMULA AND CIRCUIT

TROTTERIZATION

To simulate time-evolving process such as adiabatic quantum process, we approximate continuous process with discrete steps.

We call the discretized approximation as **Trotter** formula.

$$\exp(-i\mathcal{H}t) \approx \prod_i^n \exp\left(-i\mathcal{H}_i \frac{t}{n}\right) \quad (1)$$

where, n is a trotter steps.

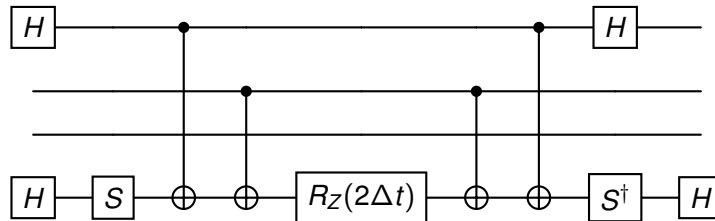
As we increase the step number n , we get more precise unitary transformation.

LIE-TROTTER FORMULA AND CIRCUIT

TROTTERIZATION

Practically, each terms of Hamiltonian are described with **Pauli string**. A single Pauli string, for example XZY , Hamiltonian has a well known corresponding circuit.

$$\exp(-i\Delta t(X \otimes Z \otimes I \otimes Y)) \quad (2)$$



OPTIMIZATION OF HAMILTONIAN

Optimization of evolution circuit is a combination of two parts.

- ▶ Mutually Commuting Partition
- ▶ Pauli-Frame

OPTIMIZATION OF HAMILTONIAN

MUTUALLY COMMUTING PARTITION

Pauli strings are always anti-commute or commute each other.

For given two Pauli strings, P_i, P_j ,

$$\text{either } [P_i, P_j] = 0 \text{ or } \{P_i, P_j\} = 0 \quad (3)$$

where, $[]$ is a commutator, and $\{ \}$ is an anti-commutator.

If all Pauli-terms of Hamiltonian are mutually commute each other, Eq(1) becomes an unitary operator of total Hamiltonian evolution of time t .

$$\exp(-i\mathcal{H}t) = \prod_i^n \exp(-i\mathcal{H}_i t) \quad (4)$$

OPTIMIZATION OF HAMILTONIAN

MUTUALLY COMMUTING PARTITION

1. We must know all commuting relation of the given Pauli-string set.
2. How to make a mutually partitions of the given set?

OPTIMIZATION OF HAMILTONIAN

MUTUALLY COMMUTING PARTITION

To make a mutually commuting partition, we have to know all commuting relationships of the given Pauli-terms of Hamiltonian. We can check the commutation with General commutativity(GC),

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If a system is n qubits system and there are m number of Pauli-terms, total operation would be, roughly,

$$\binom{m}{2} * n = O(m^2 n) \quad (5)$$

Unfortunately, $\max(m) = 2^n$ for n -qubit system Hamiltonian, it could be exponentially growth.

OPTIMIZATION OF HAMILTONIAN

MUTUALLY COMMUTING PARTITION

Chapuis et al., 2018 suggested acceleration of commuting term determination. They decompose single Pauli-string into X and Z families.

- ▶ X -family: $IIIX, XIXI, IIXI, XXII, IXXX, \dots$
- ▶ Z -family: $IIIZ, ZIZI, IIZI, ZZII, IZZZ, \dots$

$$YZIX = XIIX \cdot ZZII = x_i \cdot z_j \quad (6)$$

$$[P_i, P_j] = [x_k z_l, x_m, z_n] = \begin{cases} 0 & \text{if } [z_l, x_m] = [x_k, z_n] \\ -2P_i P_j & \text{otherwise} \end{cases} \quad (7)$$

OPTIMIZATION OF HAMILTONIAN

MUTUALLY COMMUTING PARTITION

Now, if we have compatible graph of Pauli-set, we can extract mutually commuting partition by solving a sequential Max-Clique problem of the commute graph.

It is well known NP-complete problem, from 21-complete problems. See Karp, 1972.

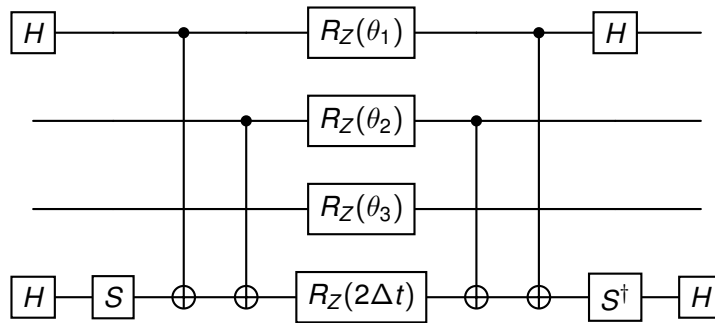
Kurita et al., 2023 suggested Ising formulation for finding Max-clique finding problem of compatible graph.

$$\mathcal{H} = -\mu_0 \sum Z_i + \mu_1 \sum h_{ij} Z_i Z_j \quad (8)$$

where, $h_{ij} = 0$ if $Z_i - Z_j$ edge weight is 0 otherwise 1, $\mu_0 = 1, \mu_1 = 2$ in Kurita et al..

OPTIMIZATION OF HAMILTONIAN

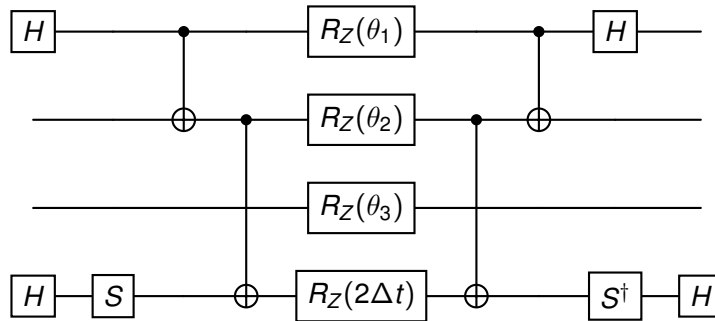
PAULI FRAME



$$\mathcal{H} = tXZIY + \theta_1XIII + \theta_2IZII + \theta_3IIII$$

OPTIMIZATION OF HAMILTONIAN

PAULI FRAME



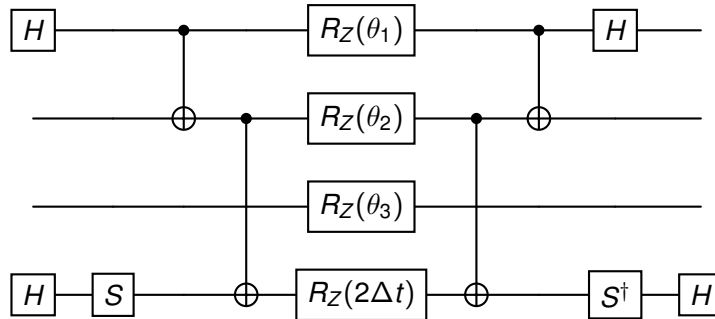
$$\mathcal{H} = tXZIY + \theta_1XIII + \theta_2XZII + \theta_3IIII$$

OPTIMIZATION OF HAMILTONIAN

PAULI FRAME

Schmitz et al., 2023 analyzed and Pauli-Frame method and optimized circuit with minimum cost of $CNOT$, H , S operations to

$$\begin{array}{ccccccc}
 Z_1 & X_1 & & X_1 & Z_1 & & X_1 & Z_1 \\
 Z_2 & X_2 & & Z_2 & Z_2 & \rightarrow_{CNOT_{1,2}} & X_1 Z_2 & Z_1 X_2 \\
 Z_3 & X_3 & \rightarrow_{H_2, H_4 - S_4} & Z_3 & X_3 & & Z_3 & X_3 \\
 Z_4 & X_4 & & Y_4 & Z_4 & & Y_4 & X_4 \\
 & & & & & \rightarrow_{CNOT_{2,4}} & X_1 Z_2 Y_4 & Z_1 X_2 Z_4
 \end{array} \tag{9}$$



$$\mathcal{H} = tXZiY + \theta_1XIII + \theta_2XZII + \theta_3IIII$$

DEGENERATE REDUCING OF MUTUALLY HAMILTONIAN

If there are two max clique on graph, sharing same number of nodes, the next Hamiltonian pick one of them randomly.

$$\mathcal{H} = -\mu_0 \sum Z_i + \mu_1 \sum h_{ij} Z_i Z_j \quad (10)$$

Eventhough, they are same in commutation graph, frame change cost can be different.

In this project, we only consider H, S costs. The weight of each Pauli-terms would be calculated with function $w(\cdot, \cdot)$, such that

- ▶ $w(\cdot, \cdot) = 0$: $(X, X), (Y, Y), (Z, Z), (Z, I)$
- ▶ $w(\cdot, \cdot) = 1$: $(X, Z), (X, Y), (X, I)$
- ▶ $w(\cdot, \cdot) = 2$: $(Y, I), (Y, Z)$

For N -qubit system, extended weight function $W(\cdot, \cdot)$ is defined as,

$$W(S_i, S_j) := \frac{1}{N} \sum_{k=1}^N w((S_i)_k, (S_j)_k) \quad (11)$$

DEGENERATE REDUCING OF MUTUALLY HAMILTONIAN

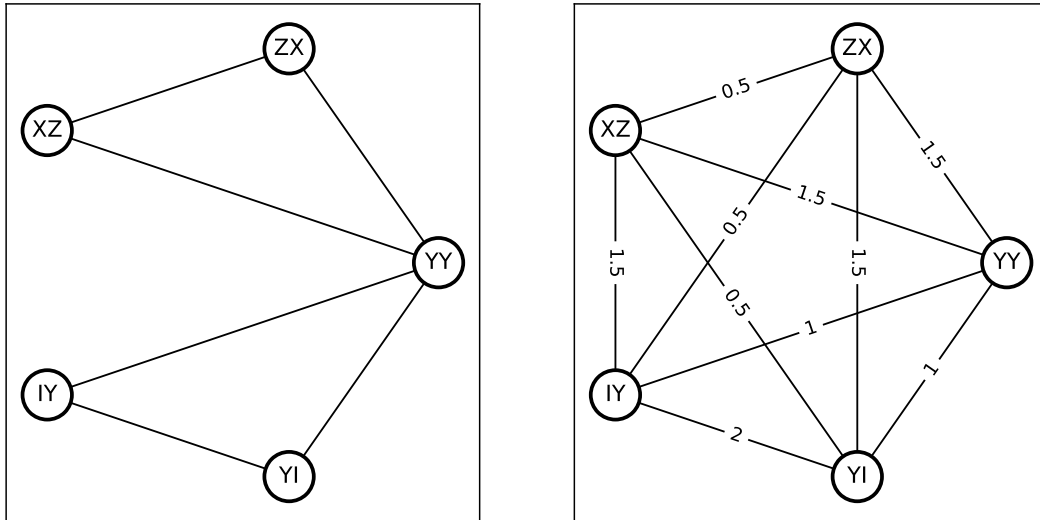


Figure. Compatible and basis transform weight graph example. Left graph is a compatible graph of 5 Pauli basis of 2 qubits system and edges are indicating commutation relationship. Right graph is a basis transform weight graph of the same Pauli-basis set of the left.

DEGENERATE REDUCING OF MUTUALLY HAMILTONIAN

We can redefine a Hamiltonian for optimization,

$$\mathcal{H} = -\mu_0 \sum Z_i + \mu_1 \sum_{i<j} h_{ij} Z_i Z_j + \mu_2 \sum_{i<j} w_{ij} Z_i Z_j \quad (12)$$

To avoid the degeneration of energy and to conserve max and commuting condition, the coefficients, μ_0, μ_1, μ_2 have next relationship.

For N qubits system,

$$\|\mu_1\| > N\|\mu_0\| \quad \|\mu_0\| > \frac{1}{2}N(N-1)\|\mu_2\| \quad (13)$$

DEGENERATE REDUCING OF MUTUALLY HAMILTONIAN

Full procedure of algorithm.

1. Find a compatible graph of the given Hamiltonian
2. Calculate weight between Pauli-strings with Eq(11)
3. Find a min-number of mutually commuting partition, p_1, p_2, \dots , using **adiabatic computer**.
4. Find a shortest hamilton path of each local partition p_i , <- reduced problem, you can use classic algorithm.
5. Connecting p_i in order to following 4 step result.

OPTIMIZATION EXAMPLE

HEH+ MOLECULAR HAMILTONIAN

Pennylane HeH+ molecule Hamiltonian:

4 qubits are required and consist of 25 Pauli-terms.

'ZXZX','IYIY','ZYZY','IZIZ','XZXZ','XIXI','YZYZ','YIYI','ZIZI', 'IIIZ', 'ZZII', 'IZZI', 'ZIIZ', 'IZII', 'IIZI', 'ZIII', 'IIZZ',
'XZXI', 'YZYI', 'XXYY', 'YXXY', 'YYXX', 'XYYX', 'IYZY', 'IXZX'

OPTIMIZATION EXAMPLE

HEH+ MOLCULAR HAMILTONIAN

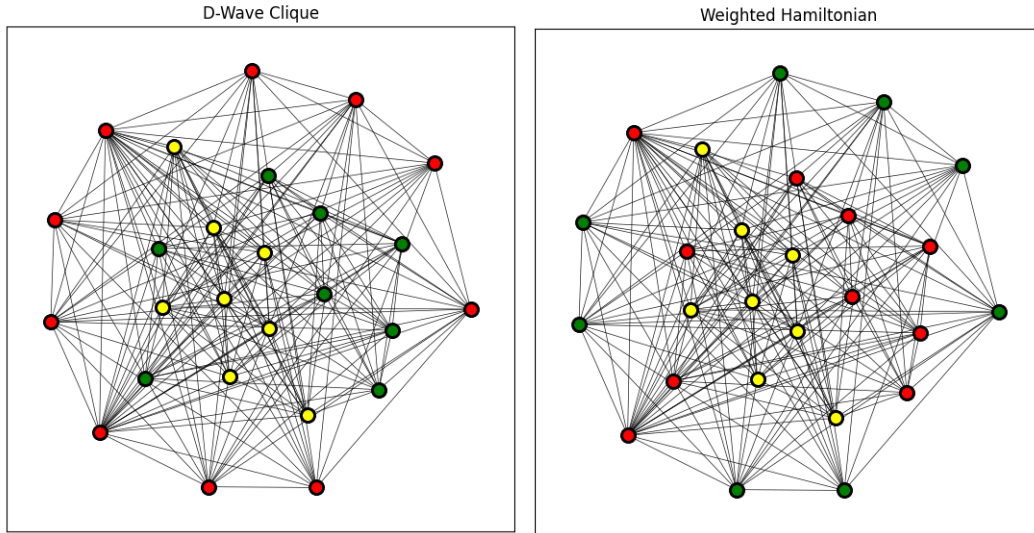


Figure. Commuting Partition HeH+ Hamiltonian Pauli-terms. Left: Ising formula solution of D-Wave. Right: Basis cost term weight added optimization.

OPTIMIZATION EXAMPLE

HEH+ MOLECULAR HAMILTONIAN

The optimization result is 3 number of partition.

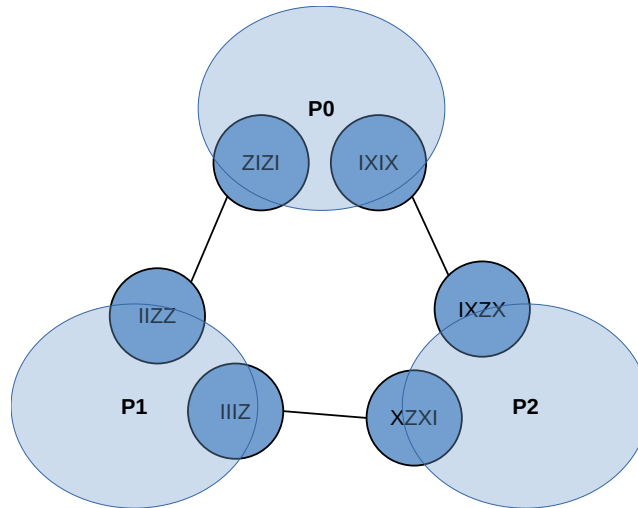
ρ_0 ['ZXZX', 'IYIY', 'ZYZY', 'IZIZ', 'XZXZ', 'XIXI', 'YZYZ', 'YIYI', 'ZIZI']

ρ_1 ['IIIZ', 'ZZII', 'IZZI', 'ZIIZ', 'IZII', 'IIZI', 'ZIII', 'IIZZ']

ρ_2 ['XZXI', 'YZYI', 'XXYY', 'YXXY', 'YYXX', 'XYYX', 'IYZY', 'IXZX']

OPTIMIZATION EXAMPLE

HEH+ MOLECULAR HAMILTONIAN



OPTIMIZATION EXAMPLE

HEH+ MOLECULAR HAMILTONIAN





Compare to PennyLane ApproxTimeEvolve() circuit

```
gates: 270
depth: 169
shots: Shots(total=None)
gate_types:
{'RZ': 106, 'CNOT': 84, 'RX': 80}
gate_sizes:
{1: 186, 2: 84}
```

```
gates: 137
depth: 107
shots: Shots(total=None)
gate_types:
{'Hadamard': 16, 'CNOT': 74, 'RZ': 25, 'S': 11, 'Adjoint(S)': 11}
gate_sizes:
{1: 63, 2: 74}
```

Figure. Left: PennyLane ApproxTimeEvolve() trotter number =1 circuit. Right: Optimized evolve circuit.

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